Query Prediction with Context Models for Populating Personal Linked Data Caches

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ABSTRACT

The emergence of a Web of Linked Data [2] enables new forms of application that require expressive query access, for which mature, Web-scale information retrieval techniques may not be suited. Rather than attempting to deliver expressive query capabilities at Web-scale, we propose the use of smaller, pre-populated data caches whose contents are personalized to the needs of an individual user. Such caches can act as personal data stores supporting a range of different applications. In this paper we formally introduce a strategy for predicting queries that can then be used to inform an a priori population of a personal cache of Linked Data harvested from Web. Based on a comprehensive user evaluation we demonstrate that our approach can accurately predict queries and their execution probability, thereby optimizing the cache population process.

Categories and Subject Descriptors

H.3.m [Information Storage and Retrieval]: Miscellaneous;
K.8.m [Personal Computing]: Miscellaneous

Keywords

Query Prediction, Context, Cache Population, Linked Data

1. INTRODUCTION

In earlier work we introduce an approach to pro-actively populate a personalized Linked Data cache for a particular, forthcoming usage context [1]. Such caches present an alternative to maintaining a centralized infrastructure that aims to serve the needs of all users in all situations. By accessing a personalized cache in the context for which it was populated, software applications may efficiently support a particular user in satisfying context-specific information needs. An important step of our cache population process is the prediction of SPARQL queries [4] that applications may execute over the cache in a given, future context. The predicted queries may then be used as a basis for the actual cache population strategy [1].

Each of the predicted queries is associated with an execution probability which represents the likelihood that the query is actually executed in the given context. Formally, we define predicted queries as follows:

Definition 1. A predicted query is a pair \((q, ep)\) where: \(q\) is a SPARQL query; \(ep \in [0, 1]\) is an estimated execution probability.

Note that the actual execution probability of predicted queries is unknown during the process of pre-populating a cache. Accordingly, the value \(ep\) associated with a predicted query \((q, ep)\) represents an assumed (i.e. predicted) probability, computed as part of the query prediction itself. Since execution probabilities indicate the relevancy of queries for our cache population approach we require a prediction strategy that forecasts the execution probability for predicted queries as accurately as possible.

The main contributions of this paper are the precise definition of our query prediction strategy and a user based evaluation thereof.

2. PRELIMINARIES

The basis of our approach is a generic context model which represents a context by a set \(C\) of context attributes (CA). Each CA is a tuple \((t, v, r)\) where \(t\) denotes the type of the CA, \(v\) is an RDF term [3] which denotes the value of the CA, and \(r \in [0, 1]\) is a relevancy score which denotes the degree to which the CA is relevant in the corresponding context. We deliberately refrain from defining a particular set of CA types for our model; a multitude of works exist that aim to capture the notion of context by defining relevant concepts (e.g. location, user interest) and their relationships. Depending on the application domain, any collection of such concepts may be suitable (and can be selected) as possible types for CAs.

In addition to a (future) context, our query prediction strategy depends on a specification of user tasks: People aim to meet their information needs by performing certain tasks. Software applications may support the performance of some of these tasks, in which case we call such a task a supported task. For example, an application may enable a user to perform the task of finding transport options to nearby places of interest.

The actual performance of supported tasks is context-specific. To capture this dependency formally we model a supported task as a pair \((P, cat)\) where \(P\) is a set of symbols that denote context dependent properties (CDP) of the task and \(cat\) is a total mapping from \(P\) to the set of all CA types in the application domain. In any possible, context-specific performance of such a task, each CDP of the task will always be instantiatied by CAs of a certain type. Mapping \(cat\) specifies this type. Formally, we define the performance of a supported task in a particular context as follows: Let \(st = (P, cat)\) be a supported task and let \(C\) be a set of CAs. A \(C\)-specific performance of \(st\) is a total, injective mapping \(perf : P \rightarrow C\) such that \(\forall cdp \in P : (perf(cdp) = (t, v, r) \Rightarrow cat(cdp) = t)\).

When a person interacts with an application in order to perform a supported task, some of the user actions will cause the application to issue queries that will be evaluated over the user’s personal cache. We assume that applications generate such queries by instantiating prepared templates for queries. To denote the set of query templates an application may instantiate when used to perform a supported task \(st = (P, cat)\) we write \(QT(st)\). Furthermore, each such query template \(q_t \in QT(st)\) is modeled as a tuple...
The likelihood that a query probability. Our estimation is based on the following assumption: C each of these templates we generate all instantiations that are pos-
plete. Hence, an additional input to our approach (in addition to
that may access the personal cache; in particular, the specified list
a complete specification of all tasks supported by all applications
may instantiate in a future context given by a set C
The general idea of our query prediction strategy is to generate all
instination of a usage scenario specification [1].

For all queries that we generate we have to estimate an execution
ment process applied to a given set of forecasted CAs. Due to space
straints we omit the definition of that process in this paper.

### 3. QUERY PREDICTION

The general idea of our query prediction strategy is to generate all
possible instantiations of those query templates that applications
may instantiate in a future context given by a set C_{lat} of CAs.
To allow for an automated implementation of this idea, we assume
a complete specification of all tasks supported by all applications
that may access the personal cache; in particular, the specified list
of query templates for these supported tasks is assumed to be com-
Hence, an additional input to our approach (in addition to
C_{lat} is a (full) set of supported tasks S_{input} which is given as part
of a usage scenario specification [1].

To predict any query that is possible w.r.t. C_{lat} and S_{input}, we
consider all query templates of all supported tasks in S_{input}. For
each of these templates we generate all instantiations that are pos-
ible in the context represented by C_{lat}. Formally, the complete set
of queries that we may generate is given as follows:

\{ \text{inst}(qt, perf) \mid st \in S_{input} \text{ and } qt \in QT(st) \}

For all queries that we generate we have to estimate an execution
probability. Our estimation is based on the following assumption:
The likelihood that a query qt = (q, ip, sub) generated for a task
st \in S_{input} is executed in the context represented by C_{lat} depends
on i) the general instantiation probability ip of qt \in QT(st) and
ii) the likelihood of perf. The first of these parameters is given
as part of ST_{input}. For computing the second parameter we intro-
duce the notion of a performance estimation method, that is, a
function that, for any finite set C of CAs and any supported task
st, maps any possible C-specific performance perf of st to a value
\xi(perf) \in [0, 1] such that \xi(perf) is an estimate for the likelihood
that st, in the context represented by C, is actually performed as
specified by perf. Examples for such functions are \xi_{min} and \xi_{avg}:

\xi_{min}(perf) = \min(\Phi(perf)) ; \xi_{avg}(perf) = \frac{1}{\Phi(perf)} \sum_{r \in \Phi(perf)} r

where

\Phi(perf) = \{ r \mid cdp \in \text{dom}(perf) \text{ and perf}(cdp) = (t, v, r) \}

Given the concept of performance estimation methods we now de-
fine the construction of a predicted query:

**Definition 2.** Let C be a set of CAs; let st = (P, cat) be a sup-
ported task; let qt = (q, ip, sub) \in QT(st) be a query template;

1In our cache population approach [1], C_{lat} results from an enrich-
ment process applied to a given set of forecasted CAs. Due to space
straints we omit the definition of that process in this paper.

let perf be a C-specific performance of st; and let \xi be a perform-
cestimation method. The \xi-based query prediction result for qt and perf, denoted by q^\xi(qt, perf), is a predicted query:

\begin{equation}
q^\xi(qt, perf) = \left( \text{inst}(qt, perf), p^{w_{ip}, w_{perf}}(ip, \xi(perf)) \right)
\end{equation}

where w_{ip}, w_{perf} \in \mathbb{N}^+ are weights and

\begin{equation}
p^{w_{ip}, w_{perf}}(ip, \xi(perf)) = (w_{ip} \cdot ip) \cdot (w_{perf} \cdot \xi(perf))
\end{equation}

Given Definition 2, the complete set of predicted queries that can
be computed with our query prediction strategy (for C_{lat}, S_{input}, and
a performance estimation method \xi) is the following:

\{ q^\xi(qt, perf) \mid st \in S_{input} \text{ and } qt \in QT(st) \}

\text{perf is a C_{lat}-specific performance of st} \}

### 4. EVALUATION

We evaluated the accuracy of our prediction of execution probabili-
ties by assessing their degree of correlation with participants’
ratings of the likelihood of executing each query; we refer to the
per query aggregate of these ratings as actual execution probabili-
ties. This was achieved by presenting participants with a concrete
description of a stranded traveller scenario and asked to rate (on
a scale of 0-10 the likelihood that, if stranded at a specific airport
with access to our hypothetical application, they would ask various
questions involving certain locations near the airport. The degree
of correlation between the predicted and actual execution probabil-
ities was calculated on a per-airport basis (and for each permutation
of weights w_{ip}, w_{perf} \in \{1, 2, 10\}) and aggregation functions \xi_{min}
and \xi_{avg} used in computing these probabilities) using Spearman’s
rank correlation coefficient \rho. Table 1 shows the highest and low-
est values of \rho for all queries across each airport. Comparison of
each \rho to the critical values (taken from [5]) in the final column
shows that all permutations of weights and aggregation functions
produced correlations that are statistically significant at the 5% al-
pha level (\alpha=0.05), for all airports. Therefore, we conclude that
our approach is able to predict, with a high degree of accuracy,
the actual execution probability of queries instantiated with a wide
range of values in the stranded traveller scenario.

### Table 1: Highest and lowest values of Spearman’s rank corre-
lation coefficient \rho for all queries across each group.

<table>
<thead>
<tr>
<th>Airport</th>
<th>N</th>
<th>Lowest \rho</th>
<th>Highest \rho</th>
<th>Crit. Val. at \alpha=0.05</th>
</tr>
</thead>
<tbody>
<tr>
<td>Coleman</td>
<td>151</td>
<td>0.216</td>
<td>0.328</td>
<td>0.165</td>
</tr>
<tr>
<td>Edmonton</td>
<td>122</td>
<td>0.289</td>
<td>0.566</td>
<td>0.165</td>
</tr>
<tr>
<td>Halifax</td>
<td>86</td>
<td>0.321</td>
<td>0.660</td>
<td>0.179</td>
</tr>
</tbody>
</table>

### 5. REFERENCES

Framework (RDF): Concepts and Abstract Syntax. W3C